OpenBSD vmm/vmd Update

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Agenda

- Where we were a year ago
- Current status
- Future plans
- Q&A

One Year Ago ...

- Reasonably complete support for OpenBSD and Linux guests
- amd64 and i386 host support
- SVM/VMX support
- Scaffolding and tools to support the above – vmd(8)/vmctl(8)

This Past Year ...

- Adding new/core features
 - Disk snapshotting
 - Template VMs
- Security Improvements
 - Removing lazy FPU support
 - L1TF mitigation
- Platform improvements
 - Bug fixing / paying down technical debt

This Past Year (cont'd) ...

- Community involvement
 - Commercial deployments of vmm hosting providers
 - Usage of vmm(4) without vmd(8) for other use cases

2018 vmm(4) Improvements

- Platform improvements
- Correctness improvements
- Performance/stability improvements
- Security improvements

• Some of these improvements impart new functionality, some are bug fixes

- Platform improvements
 - Instruction emulation improved
 - Support added for qemu fw_cfg interface
 - Support guest OS %drX registers
 - Platform support for PXE boot
 - Implement missing PIC functionality

- Instruction emulation fixes/improvement
 - RDTSCP Incorrect implementation broke SmartOS boot
 - MONITOR/MONITORX Broke booting Linux on Ryzen hosts
- QEMU fw_cfg interface support
 - Allows passing boot parameters from SeaBIOS into the VM

- Support for guest %drX registers
 - Allows hardware breakpoint usage inside guest VM
 - OpenBSD doesn't use these itself, was a subject of a security vulnerability affecting other OSes last year)

- Platform support for PXE boot
 - Implemented after last EuroBSDcon
 - Requires iPXE extension ROM image
 - Can be handled for OpenBSD guests differently (discussed later)
- Implemented missing PIC functionality
 - Basically bug fixes

2018 vmm(4) Correctness Improvements

- Correctness improvements
 - Many fixes in CPUID emulation
 - Add support for older CPUs without XSAVE
 - Handle certain SMM-related MSRs properly

2018 vmm(4) Correctness Improvements

- CPUID improvements
 - Handle misreported large leaf function #s
 - Proper topology reporting
 - Handle bizarre "rex extended CPUID" instruction used in TempleOS
 - Properly report physical address limits for the host CPU
 - Allows VMs with much larger memory

2018 vmm(4) Correctness Improvements

- Support CPUs without XSAVE
 - Older CPUs don't have this

Handle reserved SMM-related MSRs

 SDM reference guide says these should #GP on use (previously ignored, or returned 0)

2018 vmm(4) Performance Improvements

- We improved the SVM situation significantly last year ...
 - Interrupt window handling was totally broken before (fixed)
 - RFLAGS.IF handling was totally broken before (fixed)
 - Each exit would lock/unlock the kernel lock up to 4 times during exit processing before (now zero)

2018 vmm(4) Performance Improvements

- #UD on VMX instructions
 - "Inspired" by a KVM bug
 - Previously, guest usermode program could crash the VM since these instructions exit before checking CPL
 - We would terminate the VM before ...
- #GP on invalid %cr0 / %cr4 bits
 - Previously terminated the guest

2018 vmm(4) Performance Improvements

- Many of these improvements replaced "terminate the guest" with functionality appropriate for the case
 - The "terminate the guest" on anything unexpected was a remnant from early development
 - We can start to relax these conditions now

 Removed lazy FPU handling as part of the larger OS-wide effort

And of course there was L1TF last August...

- L1TF primer
 - Allows read of data in L1 cache
 - EPT addresses are treated as physical addresses (!)
 - Basically means a guest can read data out of L1 that likely was placed there while running in VMX root mode

- L1TF entry semantics (now)
 - Flush L1 cache
 - Enter guest

- How do you flush L1?
 - And is it only L1D or is there L1I \rightarrow L1D leakage too?

New microcode has "flush L1" command MSR

- What if you don't have the new microcode?
 - Read a bunch of junk, hopefully fill all of L1D what you read
 - What about the cachelines you touch after that, but before the entry (guest CPU registers)?
 - And what about L1I, anyway?

- Our L1TF 'junk' data consists of 64KB of '0xcc', just in case there is L1D → L1I leakage
 - Of course nobody who knows has said anything

 Maxime from NetBSD also reported a bug in our handling of xsetbv arguments

Thanks Maxime!

 Most of the more impactful improvements came in vmd(8) and vmctl(8)

- Qcow2 disk support
- Disk snapshots
- Template VMs
- More user friendly vmctl(8) options

2018 vmctl(8)/vmd(8) Improvements

- Qcow2 disk support
 - Supported in "standalone" or "base + snapshot" mode
 - Integrated into vmctl(8) and vmd(8)
- Old "raw" format still supported
 - Both modes "sparse" but qcow2 is "lazy allocated" (image grows over time)

Qcow2 (cont'd)

- vmctl(8) can create qcow2 disks:

-kadath- ~> vmctl create foo.raw -s 10g vmctl: raw imagefile created -kadath- ~> vmctl create foo.qcow2 -s 10g vmctl: qcow2 imagefile created -kadath- ~> ls -la foo.* -rw------ 1 mlarkin wheel 262144 Mar 18 21:30 foo.qcow2 -rw------ 1 mlarkin wheel 10737418240 Mar 18 21:30 foo.raw

Qcow2 (cont'd)

- vmctl(8) can convert disks:

-kadath- ~> vmctl create foo2.raw -i foo.qcow2						
vmctl: raw imagefile created						
-kadath- ~> ls -la foo*						
-rw	1 mlarkin	wheel	262144	Mar 18	21:30	foo.qcow2
-rw	1 mlarkin	wheel	10737418240	Mar 18	21:30	foo.raw
-rw	1 mlarkin	wheel	10737418240	Mar 18	21:33	foo2.raw

Qcow2 (cont'd)

- Sparseness is preserved:

-kadath- ~> du -h foo* 192K foo.qcow2 192K foo.raw 192K foo2.raw

Qcow2 (cont'd)

- Base image + snapshot:

-kadath- ~> vmctl create derived.qcow2 -s 10G -b foo.qcow2
vmctl: qcow2 imagefile created
-kadath- ~> ls -la *qcow2
-rw----- 1 mlarkin wheel 262144 Mar 18 21:37 derived.qcow2
-rw----- 1 mlarkin wheel 262144 Mar 18 21:30 foo.qcow2

- Qcow2 (cont'd)
 - Base image + snapshot accumulates all disk changes in snapshot disk
 - Rollback?
 - rm derived.qcow2
 - Restore previous derived.qcow2, restart VM
 - It would be nice to have rollback/rollforward be a new vmctl option (any takers?)

 vmctl(8) new command options for easier VM management

- vmctl start -B xxx
 - Set boot device (OpenBSD guests)
 - Used for autoinstalling guest VMs via network (vmctl start -B net ...)
- vmctl stop -a
 - Stop all VMs (used for shutdown scripts)

 vmctl(8) new command options for easier VM management

- vmctl stop -f

- Force kill (terminate) a VM
- Don't wait for vmmci(4)

- Template VMs
 - vmctl start -t
 - Allows for quick and easy "cloning" of VM settings

-t name Use an existing VM with the specified name as a template to create a new VM instance. The instance will inherit settings from the parent VM, except for exclusive options such as disk, interface lladdr, or interface names.

2018 vmm(4)/vmd(8) Misc Improvements

- We finally retired i386 hosts
 - It served its purpose during early development
 - Found a lot of bugs
 - Wasn't really worth maintaining anymore
- Of course i386 guests still work

2019 Goals

- We did pretty well reducing the bug count in 2018
 - But there are still many
- Solicit community involvement
 - Glad to have lots of new faces at the vmm table
- SMP is likely my personal #1 goal
 - We've done just about everything else interesting

- Underjack update
- Nested virtualization update

- Last year I talked about the underjack approach
 - Putting vmm(4) underneath the host
 - Run host as a VM itself
 - Allows XO (execute only) memory in the host
- XO memory is one defence against ROP attacks
 - Go see Todd Mortimer's talk about RETGUARD this week for another defence!

- Underjack (cont'd)
 - Kernel is working (was completed after last year's BhyveCon)
 - How do you handle running VMs in vmm(4) when the host machine itself is a VM?

- Host/root partition approach
 - Host treated as VM until launching a new (child) VM in vmm(4) via vmctl(8)
 - Temporarily exit host VM
 - Enter guest context as usual
 - Re-enter host VM context after exit
 - Repeat ad nauseum
- This approach treats the host and guest VMs as peers of each other
 - Difficult to support nested XO memory

- Nested VMX approach
 - Never leave VMX mode
 - Host VM launches VMs of its own
 - Host VM becomes nested hypervisor
 - Can more easily accomplish nested XO
- The first approach is easier to code
- The second approach allows for arbitrary levels of nesting

- Nested VMX approach status
 - Does "emulated" VMCS (no VMCS shadowing)
 - Slow

- May decide at some point to switch approaches
 - Security improvement (XO memory) vs functionality (nested VMs) decision

- Nested VMX update
 - Boots OpenBSD/vmm(4) and Linux/KVM guests
 - Needs to be redone to use shadow VMCS
 - Tons of VMCS traffic
 - Lots of issues for 32 bit hypervisor hosts if not done (HI/LO VMCS fields handled separately)
 - Maybe we don't care

- pvclock(4)
 - Paravirtualized clock
 - Modeled after KVM's PV clock interface
 - Should hopefully help time skews and high CPU usage for applications doing lots of gettimeofday() or equivalent

Community Involvement

• I'd like to take a few minutes to point out a few things going on in the community ...

- OpenBSD.amsterdam
 - Hosted vmm(4) VMs
 - Part of the hosting fee is donated to the OpenBSD foundation

Community Involvement

- OpenBSD.amsterdam (cont'd)
 - 238 VMs deployed since last year, across 7 servers

BhyveCon referral/discount code
 'BhyveCon' (5 EUR discount)

Community Involvement

- Solo5
 - Sandboxed environment for running unikernels
 - Support added for using vmm(4) as a backend hypervisor
- Would love to see more integrations like this

Questions?

Any questions?

Thank You

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